

# Scalable FileChangeNotify

## SambaXP 2014 Göttingen

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- ▶ SLA based support for more than 650 customers
  - ▶ firewalls, VPN, certificates, audits
  - ▶ based on open standards wherever possible
- ▶ Support for many OS: Linux, Cisco IOS, Windows etc.
- ▶ Compliant with BSI Grundschutz and ISO 27001 and other international regulations
- ▶ SerNet and Samba
  - ▶ technological leadership of SerNet worldwide
  - ▶ involved in almost every big European Samba project
    - ▶ SerNet distributes up-to-date Samba packages
  - ▶ samba eXPerience
    - ▶ The international Samba conference

# What is FileChangeNotify?

- ▶ MSDN on "Obtaining Directory Change Notifications":
  - ▶ An application can monitor the contents of a directory and its subdirectories by using change notifications.
- ▶ Client queries a directory handle for changes
- ▶ Filters are sent for just specific events:
  - ▶ "I'm only interested in new and deleted files"
  - ▶ "Please tell me when a file size changes"
  - ▶ ...
- ▶ API parameter **bWatchSubtree**:
  - ▶ If this parameter is TRUE, the function monitors the directory tree rooted at the specified directory.

# [MS-FSA] specification

- ▶ 2.1.1.1 Attributes per Volume (i.e. filesystem)
  - ▶ ChangeNotifyList: A list of zero or more ChangeNotifyEntries describing outstanding change notify requests for the volume.
- ▶ 2.1.4.1 Algorithm for Reporting a Change Notification for a Directory
  - ▶ **For each** ChangeNotifyEntry in Volume.ChangeNotifyList:
  - ▶ Do something like apply filters, send notifies
- ▶ "2.1.4.1," mentioned at least 12 times in [MS-FSA]
- ▶ For every metadata operation the spec makes us walk a large array
- ▶ What happens if you have 10.000 clients with 10 notifies each?
- ▶ How can we maintain the ChangeNotifyList in a cluster?

## FileChangeNotify on the wire

- ▶ Client opens a directory
- ▶ Client sends a `CHANGE_NOTIFY` request
  - ▶ FileID references the open directory handle
  - ▶ CompletionFilter shows which changes the client wants to see
  - ▶ This creates the "ChangeNotifyEntry"
- ▶ When changes happen, server replies to the `CHANGE_NOTIFY` request
- ▶ Until client sends a fresh `CHANGE_NOTIFY` request, server has to queue changes
- ▶ If the queue overflows, server can reply with "Something changed, but I don't know what"
- ▶ The ChangeNotifyEntry is only removed when closing the directory

# FileChangeNotify in Samba

- ▶ Three implementations
  - ▶ It seems that Samba often requires a few rounds to get complicated things right
  - ▶ Anyone remember the Samba 2.0 oplock implementation? :-)
- ▶ Samba 3.0
  - ▶ No global ChangeNotifyList equivalent
  - ▶ Timeout-based polling of directories per smbd
- ▶ Tridge's Samba4 implementation
  - ▶ Tridge figured out how much more of the protocol
  - ▶ Messaging-based notification
  - ▶ Ported to Samba 3.2
- ▶ Samba 4.0 notify\_index.tdb
  - ▶ Starts to make notify possible in a cluster

# Samba 3.0

- ▶ Contents of the directory are hashed
- ▶ Periodically `hash_check_notify` is called
- ▶ Recalculates the hash
- ▶ Upon changes, Samba returns `STATUS_NOTIFY_ENUM_DIR`
  - ▶ Samba did not return exactly what changed
- ▶ High load due to polling in every `smbd`
- ▶ Updates can lag
- ▶ No recursive notifies

## Samba 3.2

- ▶ During the NTVFS effort, Tridge figured out the ChangeNotifyList
- ▶ PIDL came around, complex data structures could be marshalled
- ▶ Tridge implemented the ChangeNotifyList as a hierarchical array of arrays
  - ▶ "This function is called a lot, and needs to be very fast. The unusual data structure and traversal is designed to be fast in the average case, even for large numbers of notifies"
- ▶ notify.tdb stores the ChangeNotifyList a.k.a. notify\_array in one record
- ▶ Every smbd has a copy, updated on every change
  - ▶ tdb\_seqnum was invented for this
  - ▶ This does not scale to thousands of smbds and notifies
- ▶ Problems in a cluster
  - ▶ No real tdb\_seqnum
  - ▶ One large record bounced back and forth like mad



# Samba 4.0

- ▶ The Samba 3.2-3.6 implementation has one tdb record for the complete ChangeNotifyList
- ▶ Every change pushes one huge record to every node and smbd
- ▶ Goal: Reduce write load on the central notify database
- ▶ Every notify event is path-based and needs to look at all the parents' ChangeNotifyEntry records
- ▶ Split up the notify\_array into records indexed by path
  - ▶ notify.tdb now has many path-indexed records
  - ▶ Every record holds a number of ChangeNotifyEntry records
  - ▶ A change notify event walks the path, looking for recursive entries
- ▶ Typically a lot of contention on just a few directories
  - ▶ Share root directories are very popular to look at

# Samba 4.0 clustered

- ▶ Write load on individual tdb records still high
- ▶ High n:m messaging load across nodes
  - ▶ Notify events inform many smbds, possibly many on the same node
- ▶ Split up notify.tdb into a cluster-wide notify\_index.tdb and a node-local notify.tdb
  - ▶ Both tdb's indexed by path
  - ▶ ChangeNotifyEntry records local in notify.tdb
- ▶ notify\_index.tdb holds just node numbers
  - ▶ Every node records itself with the path if any notify.tdb record exists
  - ▶ Just one single entry per node in notify\_index.tdb
- ▶ Notify events are sent to a remote proxy process
  - ▶ Proxy multi-casts notify events from its notify.tdb
- ▶ notify\_index.tdb deletion is deferred
  - ▶ Write load on notify\_index.tdb is significantly reduced
- ▶ Next bottleneck: read access on entry for "/" in notify\_index.tdb

# FileChangeNotify NextGeneration

- ▶ "This function is called a lot . . ."
  - ▶ This function (notify\_trigger) is now  $O(n)$  in the number of path components
  - ▶ For a 10-level deep file create, tdb\_parse\_record is called 10 times
  - ▶ tdb is fast, but it does cost, in particular with fcntl locks being a node-wide spinlock
- ▶ Notify events must be as cheap as possible
  - ▶ FileChangeNotify is asynchronous
  - ▶ Why not delegate notify\_trigger to some other process?
- ▶ Until a few days ago, Samba internal messaging was heavy-weight
  - ▶ tdb-based with SIGUSR1 as the async notification
  - ▶ With unix datagram messaging, sending a message is a single syscall

# Notifyd design

- ▶ Keep the ChangeNotifyList in one daemon
- ▶ Smbd adds and removes ChangeNotifyEntries by messages to notifyd
- ▶ Notify events are another type of message
- ▶ All recursive filtering is done by notifyd
- ▶ notifyd in a cluster distributes the local ChangeNotifyList

## But what about delayed messages?

- ▶ A delayed ChangeNotifyEntry creation will lose notifies
  - ▶ The notify event happens before the Entry is created
- ▶ Every message carries a timestamp
  - ▶ We could save notify events for a while
  - ▶ When should we drop them?
- ▶ Calculate a hash of the path name
  - ▶ Maintain an array of timestamps indexed by that hash
  - ▶ When an Entry comes in, check the timestamp
  - ▶ If it's later, just reply with overflow
    - ▶ All that can happen is false positives

# Questions?

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