CTDB Stories

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CTDB Project

Motivation: Support for clustered Samba

- Multiple nodes active simultaneously
- Communication between nodes (heartbeat, failover)
- Share databases between nodes

Features:

- Volatile and Persistent databases
- IP failover and load balancing
- Service monitoring

Community:

- http://ctdb.samba.org
- git://git.samba.org/ctdb.git, git://git.samba.org/samba.git

Headlines

- Merging CTDB tree in Samba tree
- Development Stories
 - High hopcount bug
 - Getting lock scheduling right
 - All nodes banned on single node failure
- Regression Stories
 - Real time or not
 - Fixing compiler warnings

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SambaXP 2013

- Merge CTDB in Samba tree?
 - Remove duplication of talloc, tdb, tevent, replace libraries
 - Autobuild testing of clustered Samba
 - Leverage off Samba release process

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CTDB tree merged with Samba

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SambaXP 2014

- To Do
 - Create waf build for CTDB, Clustered Samba
 - · Setting up clustered samba instance for autobuild
 - Split monolithic code

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June 2014

CTDB standalone waf build commited.

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Martin takes over

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- Replace dependency on ctdb-util with samba-util
- Hook CTDB into top level using --with-cluster-support

November 2014

CTDB build integrated into toplevel build.

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CTDB Releases

- 2.5.4 (September 2014) 156 patches
 - Support for TDB robust mutexes
 - Add ctdb detach
 - Avoid running ctdb helpers at real-time priority
 - Improved vacuuming performance
- 2.5.5 (April 2015) 119 patches
 - Fix handling of IPv6 addresses
 - Fix regression in socket handling code
 - Make statd-callout scalable

Contributions in 2014

- 196 Martin Schwenke
- 184 Amitay Isaacs
 - 55 Michael Adam
 - 10 Volker Lendecke
 - 3 Srikrishan Malik
 - 3 Andrew Bartlett
 - 2 Stefan Metzmacher
 - 2 Gregor Beck
 - 2 Bjorn Baumbach
 - 1 Matthias Dieter Wallnofer
 - 1 Jeremy Allison
 - 1 Ira Cooper
 - 1 David Disseldorp

Contributions since Jan 2015

- 118 Martin Schwenke
 - 15 Amitay Isaacs
 - 12 Volker Lendecke
 - 3 Rajesh Joseph
 - 1 Michael Adam
 - 1 Led
 - 1 Jelmer Vernooij
 - 1 David Disseldorp
 - 1 Christof Schmitt

High hopcount bug

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Problem

Logs filled with entries like:

ctdbd: High hopcount 2823099 dbid:0x7a19d84d key:0x6f9f65c4

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```
static void ctdb_call_send_redirect(ctdb, ctdb_db, key, c, header)
{
    uint32_t lmaster = ctdb_lmaster(ctdb, &key);
    c->hdr.destnode = lmaster;
    if (ctdb->pnn == lmaster) {
        c->hdr.destnode = header->dmaster:
    3
    c->hopcount++;
    if (c->hopcount%100 > 95) {
        DEBUG(DEBUG_WARNING,("High hopcount ..."));
    3
    ctdb_queue_packet(ctdb, &c->hdr);
```


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- Reply to Node 0 (DMASTER_REPLY)



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- Reply to Node 1 (DMASTER_REQ)
- Reply to Node 0 (DMASTER_REPLY)
- Reply to Client (REPLY_CALL)

Debugging

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Suspects

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- Record header corruption
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 - However, the problem did not go away
- Locking code was being modified

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UPDATE	db[notify_index.tdb]:	store:	hash[0x0aa13d47]	rsn[9620]	dmaster[1]
UPDATE	db[notify_index.tdb]:	store:	hash[0x0aa13d47]	rsn[9621]	dmaster[1]
UPDATE	db[notify_index.tdb]:	store:	hash[0x0aa13d47]	rsn[9622]	dmaster[1]
UPDATE	db[notify_index.tdb]:	store:	hash[0x0aa13d47]	rsn[9623]	dmaster[1]

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Node 0 requests the record. Node 1 updates DMASTER.

UPDATE db[notify_index.tdb]: store: hash[0x0aa13d47] rsn[9640] dmaster[1] UPDATE db[notify_index.tdb]: store: hash[0x0aa13d47] rsn[9641] dmaster[1] UPDATE db[notify_index.tdb]: store: hash[0x0aa13d47] rsn[9641] dmaster[0]

And Node 1 migrates the record to Node 0

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 - Creates a lock request

• Meanwhile, more record requests queue up

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Waiting reqid:732 key:0x0aa13d47 Waiting reqid:684 key:0x0aa13d47 Waiting reqid:715 key:0x0aa13d47 Waiting reqid:701 key:0x0aa13d47 • Meanwhile, more record requests queue up

```
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Waiting reqid:684 key:0x0aa13d47
Waiting reqid:715 key:0x0aa13d47
Waiting reqid:701 key:0x0aa13d47
```

• Soon after high hopcount messages are logged on Node 0

High hopcount 97 key:0x0aa13d47 reqid=00004771 pnn:0 src:1 lmaster:1 High hopcount 99 key:0x0aa13d47 reqid=00004771 pnn:0 src:1 lmaster:1 High hopcount 196 key:0x0aa13d47 reqid=000039f9 pnn:0 src:0 lmaster:1 High hopcount 198 key:0x0aa13d47 reqid=000039f9 pnn:0 src:0 lmaster:1

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Waiting reqid:732 key:0x0aa13d47
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• These record requests bounce very quickly. After 2 seconds:

High hopcount 955596 key:0x0aa13d47 reqid=000039f9 pnn:0 src:0 lmaster:1 High hopcount 955598 key:0x0aa13d47 reqid=000039f9 pnn:0 src:0 lmaster:1 High hopcount 955597 key:0x0aa13d47 reqid=00004771 pnn:0 src:1 lmaster:1 High hopcount 955599 key:0x0aa13d47 reqid=00004771 pnn:0 src:1 lmaster:1

• Sometime later the migrated record request gets processed

• And the bouncing requests stop.

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Solution

Avoid processing record requests for record in migration
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Why lock scheduling

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- There are multiple databases
- Freeze requests are handled independently

• New locking API abstaction - Naive approach

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... till database recovery is triggered under load

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- Maximum number of active lock requests
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... till database recovery is triggered under load

- Active queue is full and freeze lock requests are pending
- Freeze lock requests need to be scheduled immediately

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Performance is not good when record locking is in use

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Solution

• A single limit on active records kills performance for locking requests across multiple databases

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- Rely on kernel to do "fair scheduling"
- Before scheduling a lock request, check if there is an active lock request for the same record

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CTDB is consuming 100% CPU under heavy load



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Solution

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Better Solution

• Use better data structure for checking active lock requests

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 - Relies on parsing /proc/locks
- Cannot be used with TDB robust mutexes
- Recreate after disabling TDB robust mutexes

• CTDB fails to freeze smbXsrv_session_global.tdb

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ctdbd-lock: /usr/bin/ctdb_lock_helper smbXsrv_session_global.tdb.0 168 223318 ctdbd-lock: /usr/bin/ctdb_lock_helper smbXsrv_tcon_global.tdb.0 168 EOF ctdbd-lock: /usr/sbin/smbd smbXsrv_tcon_global.tdb.0 251880 251880 W ctdbd-lock: /usr/bin/ctdb_lock_helper locking.tdb.0 168 EOF ctdbd-lock: /usr/bin/ctdb_lock_helper smbXsrv_open_global.tdb.0 168 EOF ctdbd-lock: /usr/bin/ctdb_lock_helper cnscm_monitoring.tdb.0 168 EOF ctdbd-lock: /usr/bin/ctdb_lock_helper cnscm_monitoring.tdb.0 168 EOF

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• Samba process is holding a lock

Stack trace for relevant samba process

#0 0x00007fde05236218 in poll () from /lib64/libc.so.6 0x00007fde0863a93c in poll_one_fd () #1 #2 0x00007fde0861146b in ctdb_packet_fd_read_sync_timeout () 0x00007fde08611c0d in ctdb_packet_fd_read_sync () #3 0x00007fde086126fa in ctdb_read_req () #4 #5 0x00007fde08612eae in ctdbd_parse () 0x00007fde0862184d in db_ctdb_parse_record () #6 #7 0x00007fde0861d9d4 in dbwrap_parse_record () #8 0x00007fde0861dc2a in dbwrap_fetch () 0x00007fde086250fd in dbwrap_watch_record_stored () #9 #10 0x00007fde0861dc86 in dbwrap_record_delete () #11 0x00007fde083887bd in smbXsrv_session_logoff () #12 0x00007fde083892aa in smbXsrv_session_logoff_all_callback () #13 0x00007fde08626389 in db_rbt_traverse_internal () #14 0x00007fde086264da in db rbt traverse () #15 0x00007fde0861d96a in dbwrap_traverse () #16 0x00007fde08389918 in smbXsrv_session_logoff_all () #17 0x00007fde088e41a0 in exit server common () #18 0x00007fde088e462e in smbd_exit_server_cleanly () #19 0x00007fde083609e2 in exit_server_cleanly ()

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Amitay Isaacs CTDB Stories

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 - Process all pending call requests for that database

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- CTDB_MANAGES_SAMBA=yes
- In 50.samba, startup event starts smbd
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- Samba not getting scheduled to read from CTDB
- If write() calls fails with EAGAIN, back off

Questions/Comments?



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